OLLSCOIL NA hÉIREANN, MÁ NUAD
NATIONAL UNIVERSITY OF IRELAND, MAYNOOTH

THIRD COMPUTER SCIENCE AND SOFTWARE ENGINEERING EXAMINATION

## SAMPLE 2001

## PAPER SE307

## COMPLEXITY THEORY

Dr. A. Mycroft, Mr. S. Brown, Mr. T. Naughton
Answer ALL QUESTIONS from Section A and any TWO questions from Section B. Negative marking will be applied for Section $A$ ( 2 marks for a correct answer, $-\mathbf{0 . 5}$ for incorrect, 0 for no attempt). Time Allowed: 2 hours.

## SECTION A (30 marks)

1. Given an alphabet $\Sigma$, a language over $\Sigma$ is a
(a) superset of $2^{\Sigma}$
(b) subset of $\Sigma^{*}$
(c) superset of $\Sigma^{*}$
(d) proper subset of $\Sigma_{0}^{*}$
(e) none of the above
2. Which of the following functions has the largest growth rate?
(a) $n^{1 / 2}$
(b) $n^{1 / 10}$
(c) $n^{100}$
(d) $2^{n / 2}$
(e) $2^{n!}$
3. When $n$ doubles
(a) $\log n$ increments
(b) a linear function in $n$ doubles
(c) a quadratic function in $n$ quadruples
(d) a quadratic function in $n$ squares
(e) an exponential function in $n$ squares
4. The 'strongest' of the following statements we can make about the summation [2 marks] $1+2+3+\ldots+n$ is
(a) $=O\left(n^{2}\right)$
(b) $=\frac{1}{2} n^{2}+O(n)$
(c) $=O\left(n^{2}\right)+O(n)$
(d) $=O\left(n^{2}\right)-n$
(e) $=n^{2}+O(n)$
5. The sum $1^{2}+2^{2}+3^{2}+\ldots+n^{2}=\frac{1}{3} n\left(n+\frac{1}{2}\right)(n+1)$ would also be
(a) $=O\left(n^{3}\right)$
(b) $=\frac{1}{3} n+O\left(n^{3}\right)$
(c) $=O\left(n^{4}\right)$
(d) all of the above
(e) none of the above
6. We can reduce our number of unique tape symbols in a TM table of behaviour by [2 marks]
(a) increasing the size of the set of symbols
(b) writing more symbols on the tape
(c) increasing the number of states of mind
(d) a method requiring both (b) and (c)
(e) either (b) or (c)
7. The minimum number of symbols we can reduce our set of symbols to without [2 marks] changing the functionality of any possible TM is
(a) 0
(b) 1
(c) 2
(d) greater than 2
(e) no limit
8. Which of the following are not one of the 'unrestricted' models of computation?
(a) TMs with only one tape
(b) $k$-tape TMs with a finite set of symbols
(c) $k$-tape TMs whose tapes are infinite in one direction only
(d) RAMs with fixed sized registers
(e) none of the above
9. Given a countable set $\mathbf{A}$ of subsets of a set $\mathbf{X}$, where each $a \in \mathbf{A} \Rightarrow a \subseteq \mathbf{X}$, it can [2 marks] be said that
(a) $\mathbf{A} \neq 2^{\mathbf{X}}$
(b) $\mathbf{A}=2^{\mathbf{X}}$
(c) $($ cardinality of $\mathbf{A})>\left(\right.$ cardinality of $2^{\mathbf{X}}$ )
(d) if $\mathbf{X}$ is not finite then $\mathbf{A}$ is not uncountable
(e) none of the above
10. Which of the following languages is not recursively enumerable?
(a) $\{\mathrm{a}, \mathrm{b}, \mathrm{c}\}$
(b) the odd integers
(c) the prime numbers
(d) the halting Turing machines
(e) none of the above
11. Which of the following languages is not recursive?
(a) $\{\mathrm{a}, \mathrm{b}, \mathrm{c}\}$
(b) the odd integers
(c) the prime numbers
(d) the halting Turing machines
(e) none of the above
12. The infinite set of all words over an alphabet $\Sigma$ is denoted
(a) $\Sigma^{*}$
(b) $\Sigma_{0}^{*}$
(c) $|\Sigma|$
(d) $2^{\Sigma}$
(e) none of the above
13. Given the diagram below, depicting sets $\mathbf{X} \subseteq \mathbf{Y} \subseteq \mathbf{Z}$ and elements $a \in \mathbf{X}, b \in \mathbf{Y}$, [2 marks] $c \in \mathbf{Z}$, which of the following statements is false?

(a) $b$ is $\mathbf{X}$-hard
(b) $b$ is $\mathbf{Y}$-complete
(c) $c$ is Z-hard
(d) $b$ is Z-hard
(e) none of the above
14. Given a countable set $\mathbf{A}$ of subsets of a set $\mathbf{X}$, where each $a \in \mathbf{A} \Rightarrow a \subseteq \mathbf{X}$, it can [2 marks] be said that
(a) $\mathbf{A} \neq 2^{\mathbf{X}}$
(b) $\mathbf{A}=2^{\mathbf{X}}$
(c) $|\mathbf{A}|>\left|2^{\mathbf{X}}\right|$
(d) if $\mathbf{X}$ is not finite then $\mathbf{A}$ is not uncountable
(e) none of the above
15. Given that a $k$-tape deterministic Turing machine $T$ with $k \geq 1$ can be defined by [2 marks] the tuple $\left\langle Q, \Sigma, I, q_{0}, F\right\rangle$ which of the following is false?
(a) $\Sigma$ always includes a 'blank' symbol
(b) $I$ is a set of quintuples
(c) $q_{0} \in Q$ is always the initial state
(d) $|F| \geq|Q|$
(e) none of the above

## SECTION B (40 marks)

1. You are given a 2-tape Turing machine $T=\left\langle Q, \Sigma, I, q_{0}, F\right\rangle=$ $\langle\{00,01,02,03,09\},\{0,1,-\}, I, 00,\{09\}\rangle$ that operates on binary strings. As usual, the head of the first tape will be positioned at the beginning of the input and the second tape will be blank. $I$ is

| $q$ | $s$ | $q^{\prime}$ | $s^{\prime}$ | $m$ |
| :---: | :---: | :---: | :---: | :---: |
| 00 | $\langle 0,-\rangle$ | 00 | $\langle 0,-\rangle$ | $\langle\mathrm{R}, \mathbf{S}\rangle$ |
| 00 | $\langle 1,-\rangle$ | 01 | $\langle 1,-\rangle$ | $\langle\mathrm{R}, \mathbf{S}\rangle$ |
| 00 | $\langle-,-\rangle$ | 02 | $\langle-,-\rangle$ | $\langle\mathrm{L}, \mathbf{S}\rangle$ |
| 01 | $\langle 0,-\rangle$ | 00 | $\langle 0,-\rangle$ | $\langle\mathrm{R}, \mathbf{S}\rangle$ |
| 01 | $\langle 1,-\rangle$ | 01 | $\langle 1,-\rangle$ | $\langle\mathrm{R}, \mathbf{S}\rangle$ |
| 01 | $\langle-,-\rangle$ | 03 | $\langle-,-\rangle$ | $\langle\mathrm{L}, \mathbf{S}\rangle$ |
| 02 | $\langle 0,-\rangle$ | 09 | $\langle 0,0\rangle$ | $\langle\mathbf{S}, \mathbf{S}\rangle$ |
| 02 | $\langle 1,-\rangle$ | 09 | $\langle 1,0\rangle$ | $\langle\mathbf{S}, \mathbf{S}\rangle$ |
| 02 | $\langle-,-\rangle$ | 09 | $\langle-,-\rangle$ | $\langle\mathrm{R}, \mathbf{S}\rangle$ |
| 03 | $\langle 0,-\rangle$ | 09 | $\langle 0,1\rangle$ | $\langle\mathbf{S}, \mathbf{S}\rangle$ |
| 03 | $\langle 1,-\rangle$ | 09 | $\langle 1,1\rangle$ | $\langle\mathbf{S}, \mathbf{S}\rangle$ |
| 03 | $\langle-,-\rangle$ | 09 | $\langle-,-\rangle$ | $\langle\mathrm{R}, \mathbf{S}\rangle$ |

(a) What does $T$ do? Give as concise an explanation as you can.
(b) Convert $T$ into a functionally-identical Turing machine that requires at least [7 marks] one less state than $T$.
(c) What would be the implications if a RAM algorithm solving an NP-complete problem was found to have (i) an exponential upper bound, or (ii) an exponential lower bound?
2. (a) Explain how we can view the construction of a deterministic Turing machine as a search through an ordered set.
(b) Outline how a RAM (random-access machine) could be simulated on a Turing machine. Explain how your simulation lends weight to the Invariance thesis.
(c) What does a polynomial reduction $A \leq B$ between two problems establish about their relative complexities? How could one use a reduction to prove non-membership of a class?
3. (a) Define the 1-D (one-dimensional) tiling problem and prove that it is decidable.
(b) At a recent Clubs and Societies awards function a particularly bored individual got to thinking about whether it would be possible to nominate $k$ clubs or societies for awards such that every participating student was a member of one, but only one, nominated club or society. Prove this problem is NP-complete.

